

# Concurrency and Privacy with Payment Channel Networks

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# Bitcoin Scalability Issues

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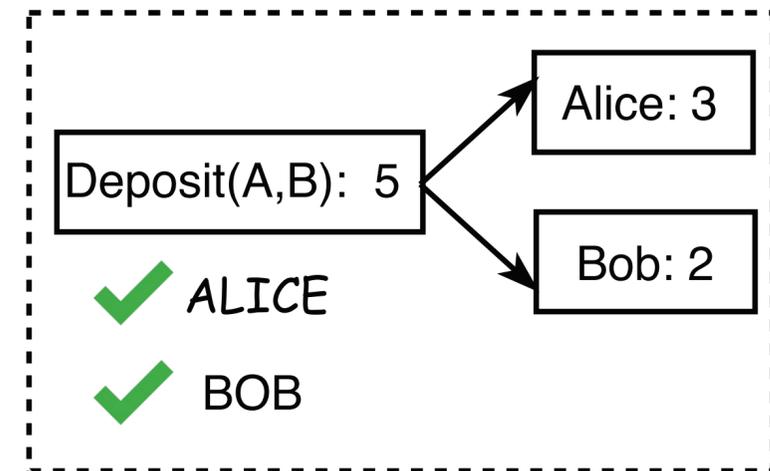
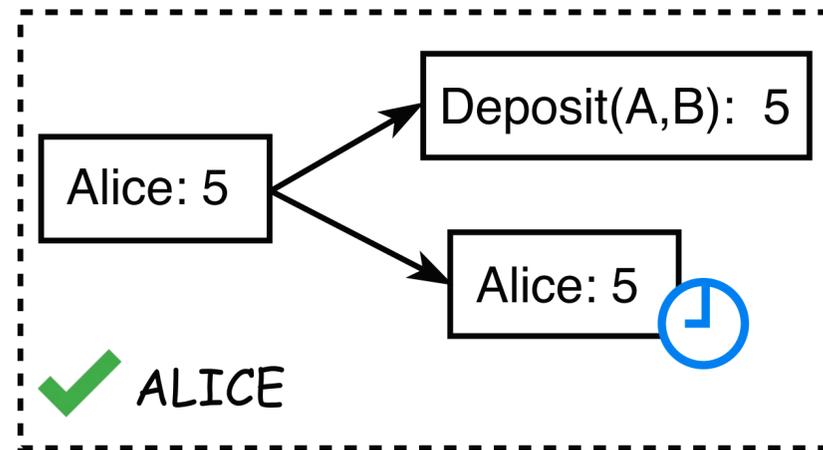
- ❖  $< 10$  transactions per second
- ❖  $> 135$  GB of memory required
- ❖ No micropayment (high fees)



# Payment Channels

- ❖ Enable multiple payments between two users without committing every single payment to the blockchain

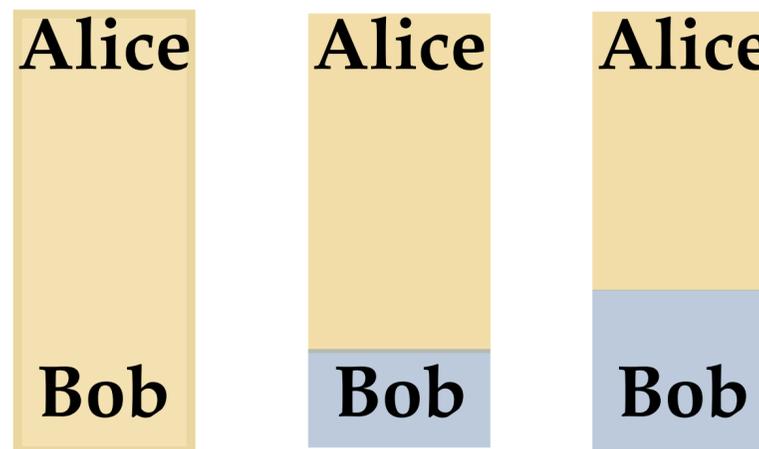
## Blockchain Transactions



OPEN CHANNEL

CLOSE CHANNEL

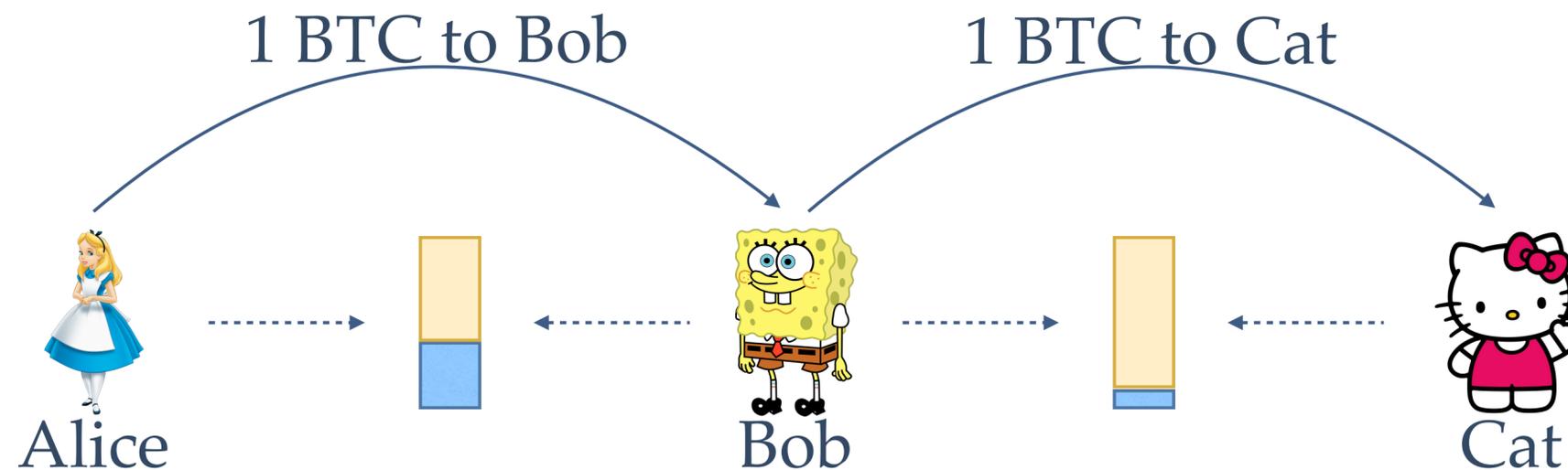
## Off-chain Payments



# Payment Channel Networks (PCN)

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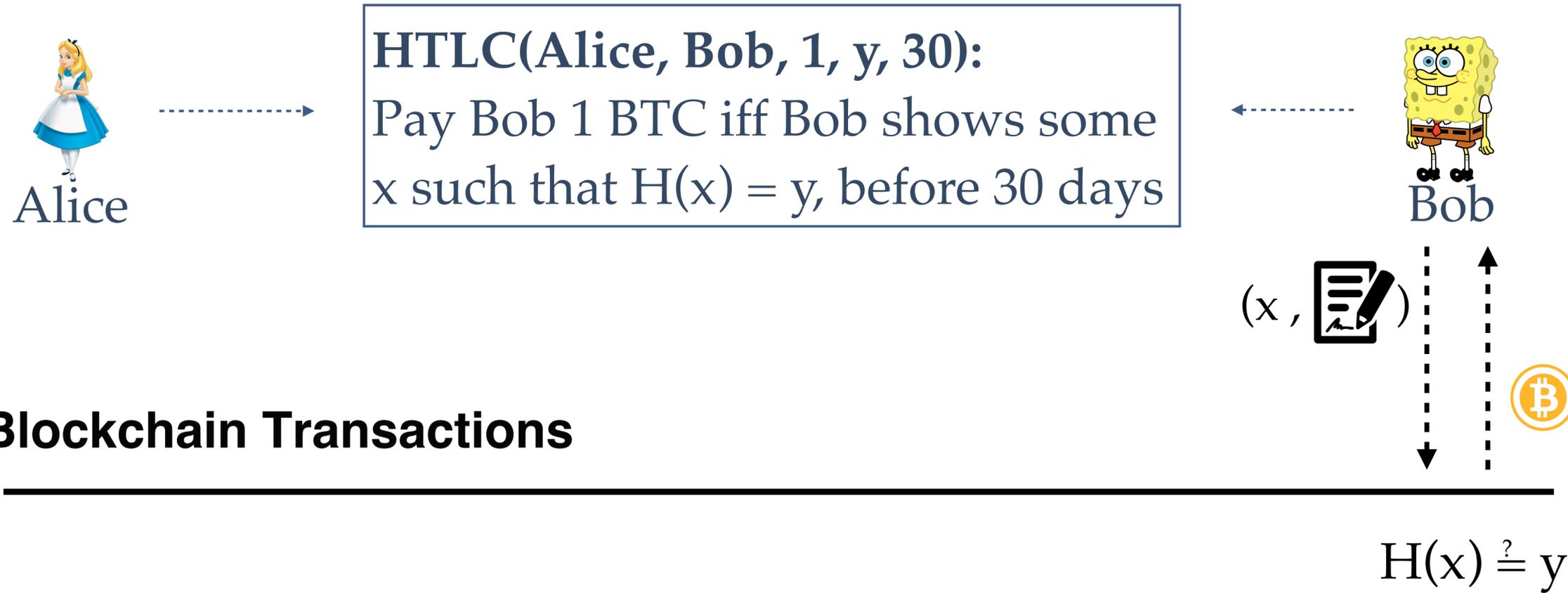
- ❖ Each payment channel requires to deposit bitcoins
- ❖ Impractical to open a channel with every other user



# Hash Time-Lock Contracts

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- ❖ Hash-Time Lock Contract (HTLC) enables conditional payments between two users

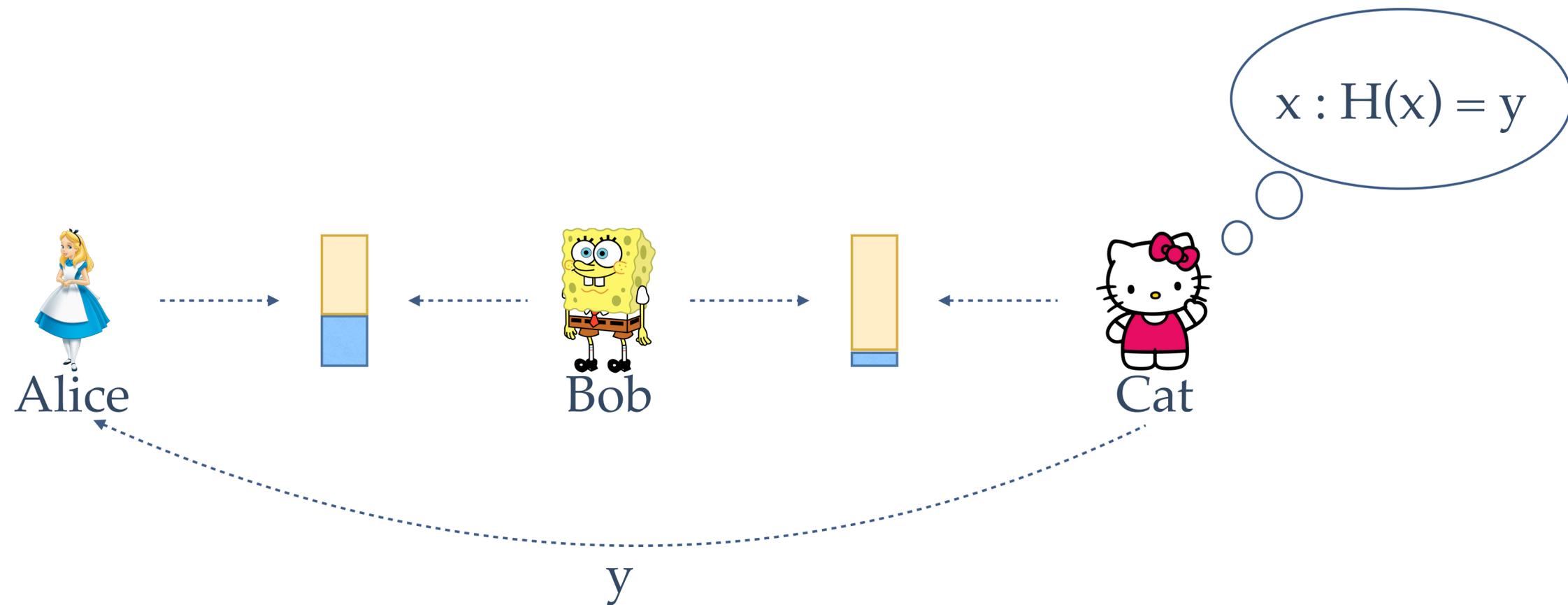


**Blockchain Transactions**

# The Lightning Network

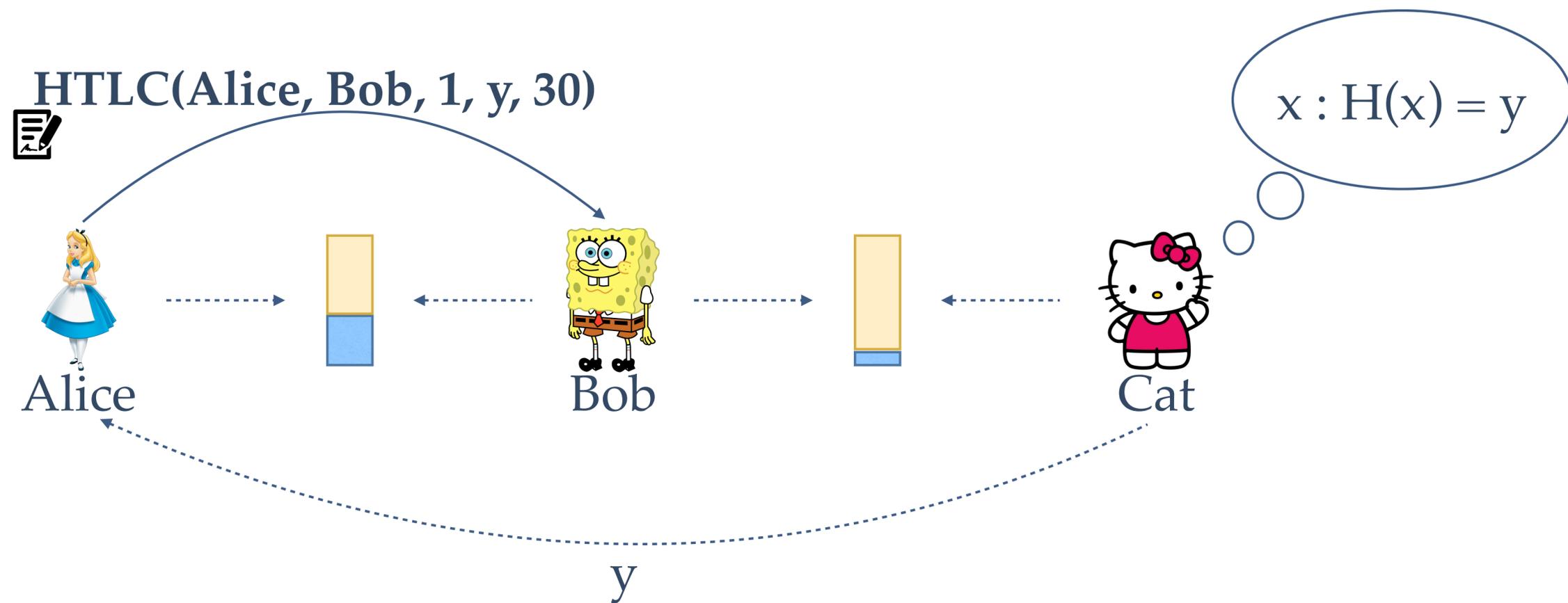
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- ❖ Multiple “chained” HTLC enable multi-hop payments in the presence of untrusted intermediaries



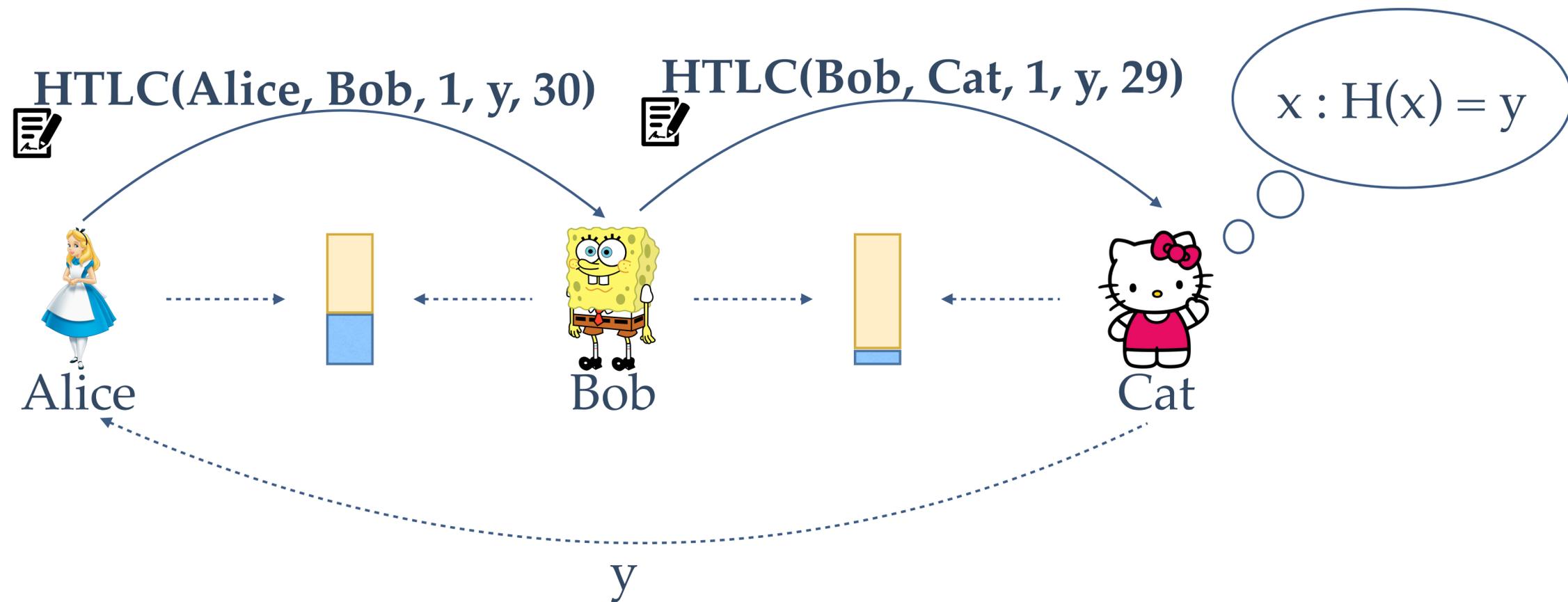
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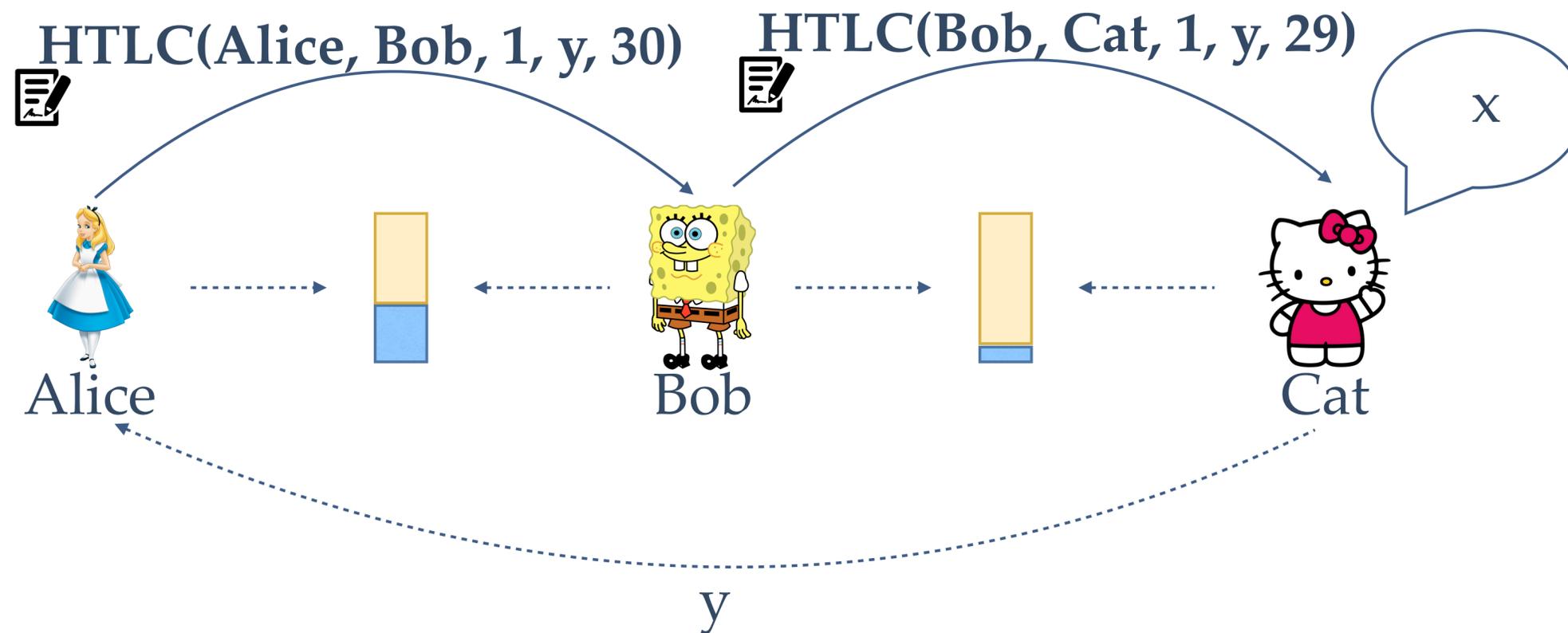
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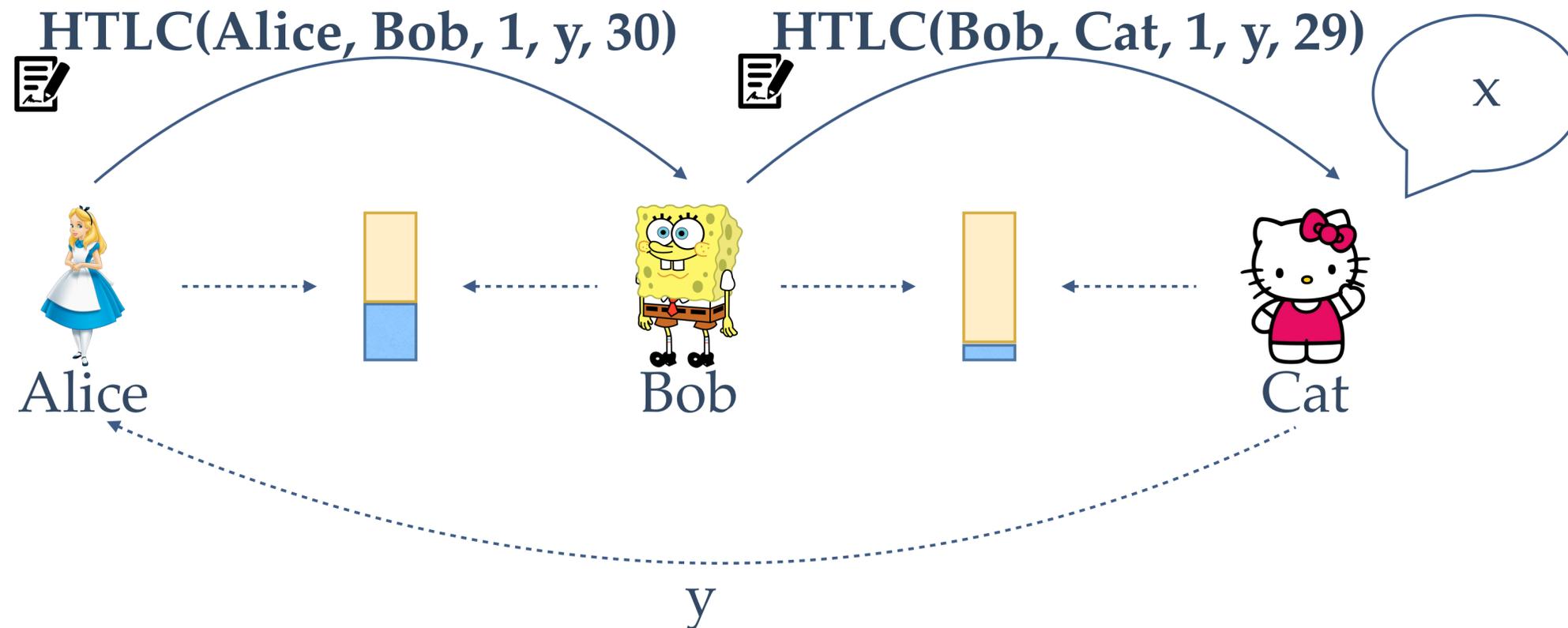
# The Lightning Network

- ❖ Multiple “chained” HTLC enable multi-hop payments in the presence of untrusted intermediaries



# The Lightning Network

- ❖ Multiple “chained” HTLC enables multi-hop payments in the presence of untrusted intermediaries
- ❖ Bob does not gain or lose coins



# Contributions

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- ❖ Definition of security and privacy properties for PCNs
- ❖ Privacy analysis of PCNs and solution (Fulgor)
- ❖ Concurrency analysis of PCNs and solution (Rayo)
- ❖ Prototype implementation

# Security Properties

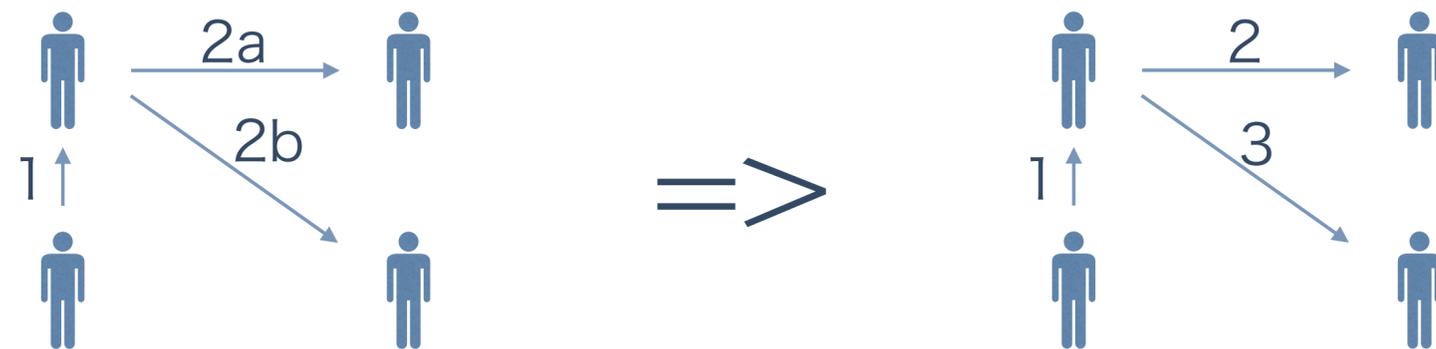
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❖ Our model highlights two main security properties:

❖ **Balance Security:**

$$\Pr \left[ \begin{array}{c} \Delta = 10 \\ \text{Red} \rightarrow \text{Red} \rightarrow \text{Blue} \rightarrow \text{Red} \rightarrow \text{Red} \end{array} \xrightarrow{\text{after payment}} \begin{array}{c} \Delta = 9 \\ \text{Red} \rightarrow \text{Red} \rightarrow \text{Blue} \rightarrow \text{Red} \rightarrow \text{Red} \end{array} \right] < \text{negl}$$

❖ **Serializability:**



# Privacy Properties

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❖ Our model highlights two privacy properties

❖ (Off-path) value privacy:



❖ (On-path) relationship anonymity:



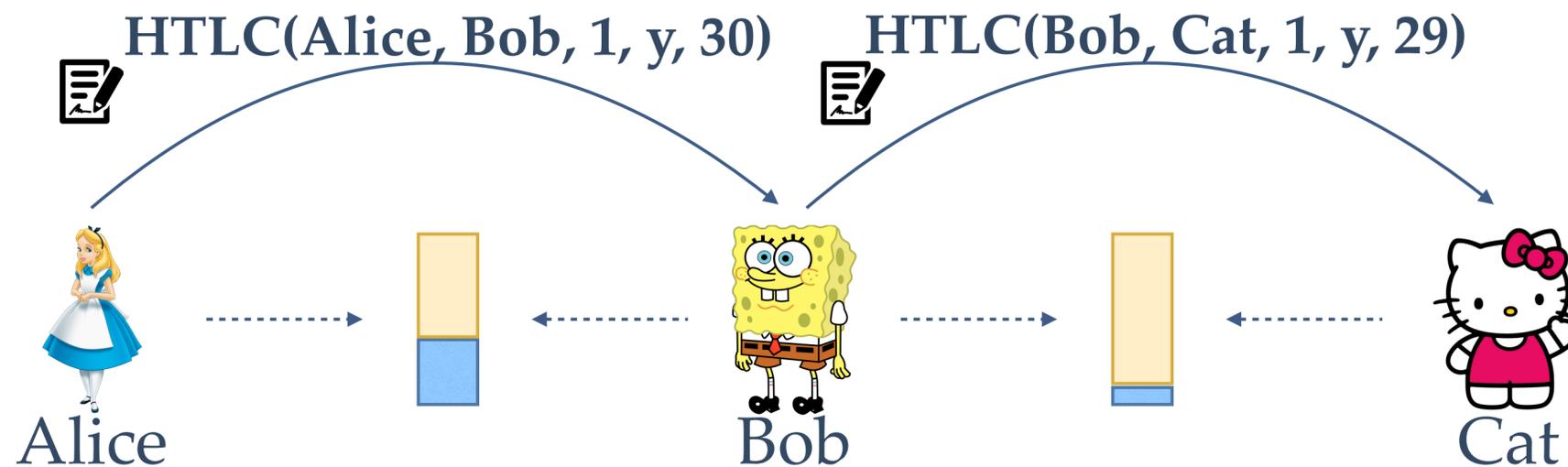
# Privacy in PCNs: Challenge?

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- ❖ Off-chain payments  $\Rightarrow$  Privacy-preserving payments

## Blockchain Transactions

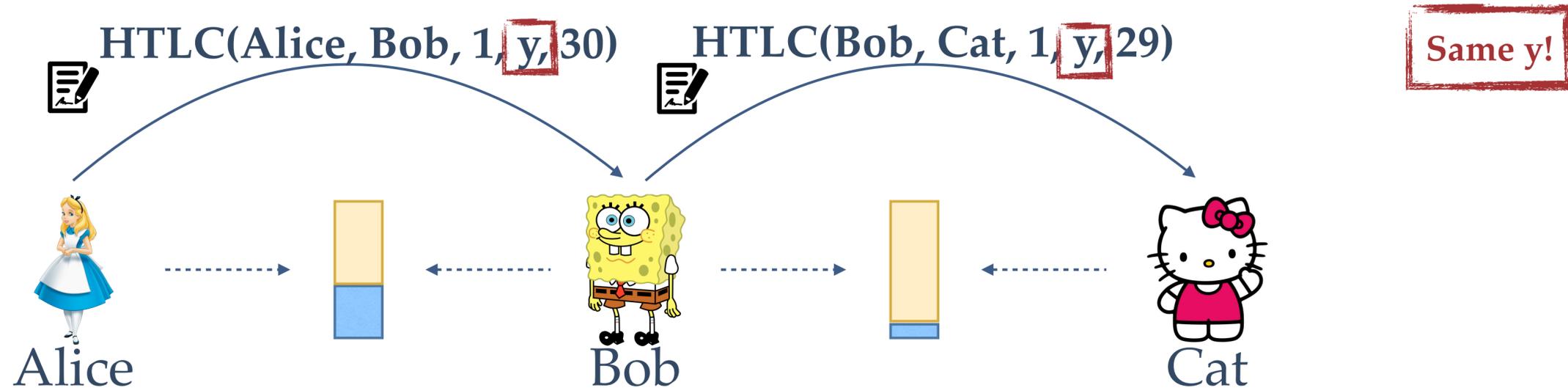
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# Privacy in PCNs: Challenge?

❖ Off-chain payments  $\neq$  Privacy-preserving payments

## Blockchain Transactions



# Privacy in PCNs: Our Solution

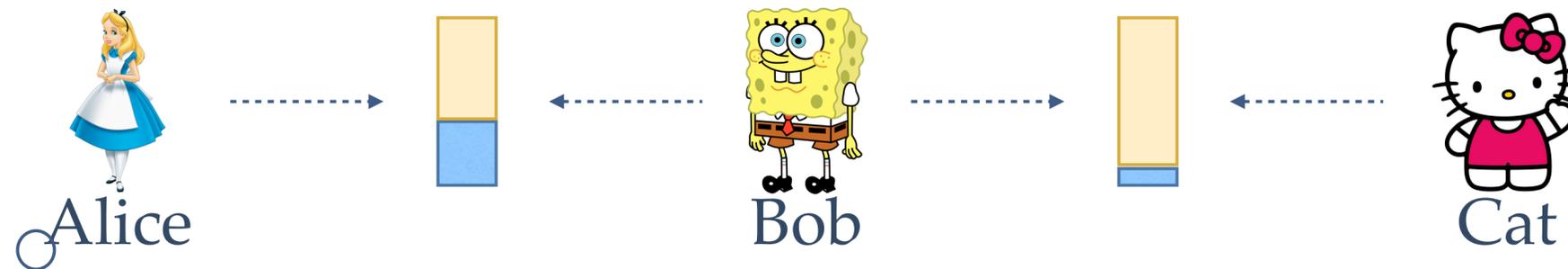
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- ❖ Our setting: P2P Network
- ❖ Our goal:
  - ❖ On-chain operations: HTLC as in the Lightning Network
  - ❖ Rest of cryptographic operations must be off-chain
  - ❖ Full compatibility with the current Bitcoin script
- ❖ Our solution:
  - ❖ Fulgor: Based on Multi-hop HTLC

# Multi-hop HTLC

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- ❖ Building block: Non-interactive zero knowledge (ZKBoo [GMO16])

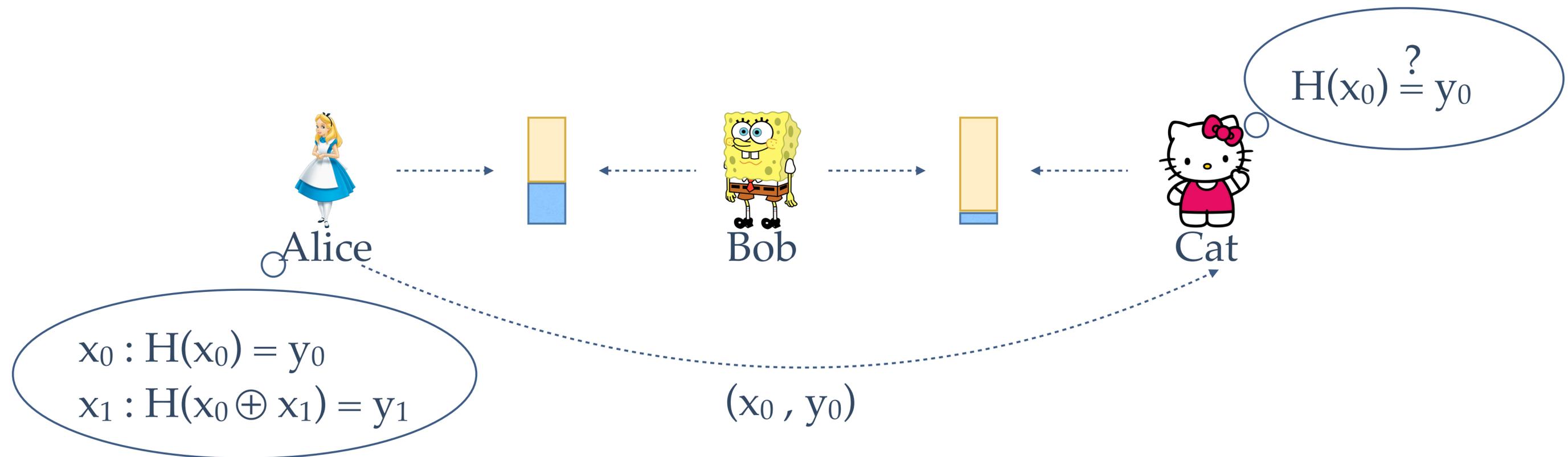


$$x_0 : H(x_0) = y_0$$

$$x_1 : H(x_0 \oplus x_1) = y_1$$

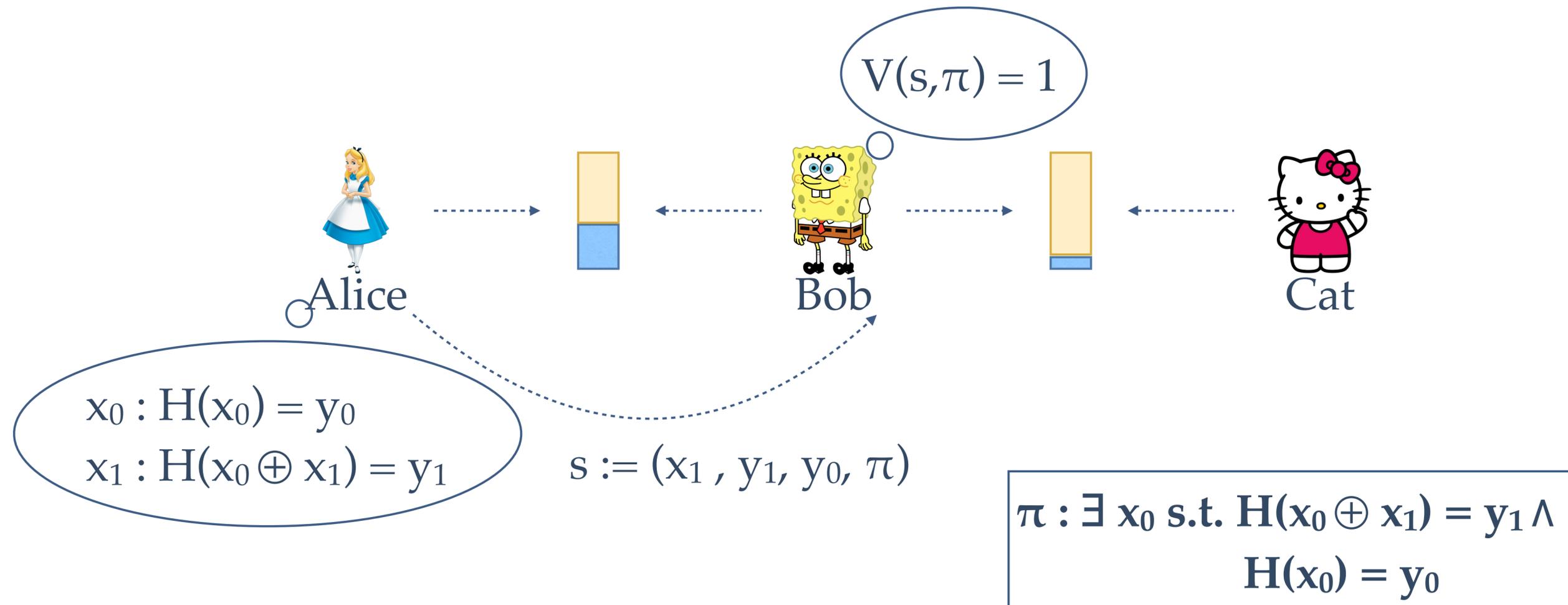
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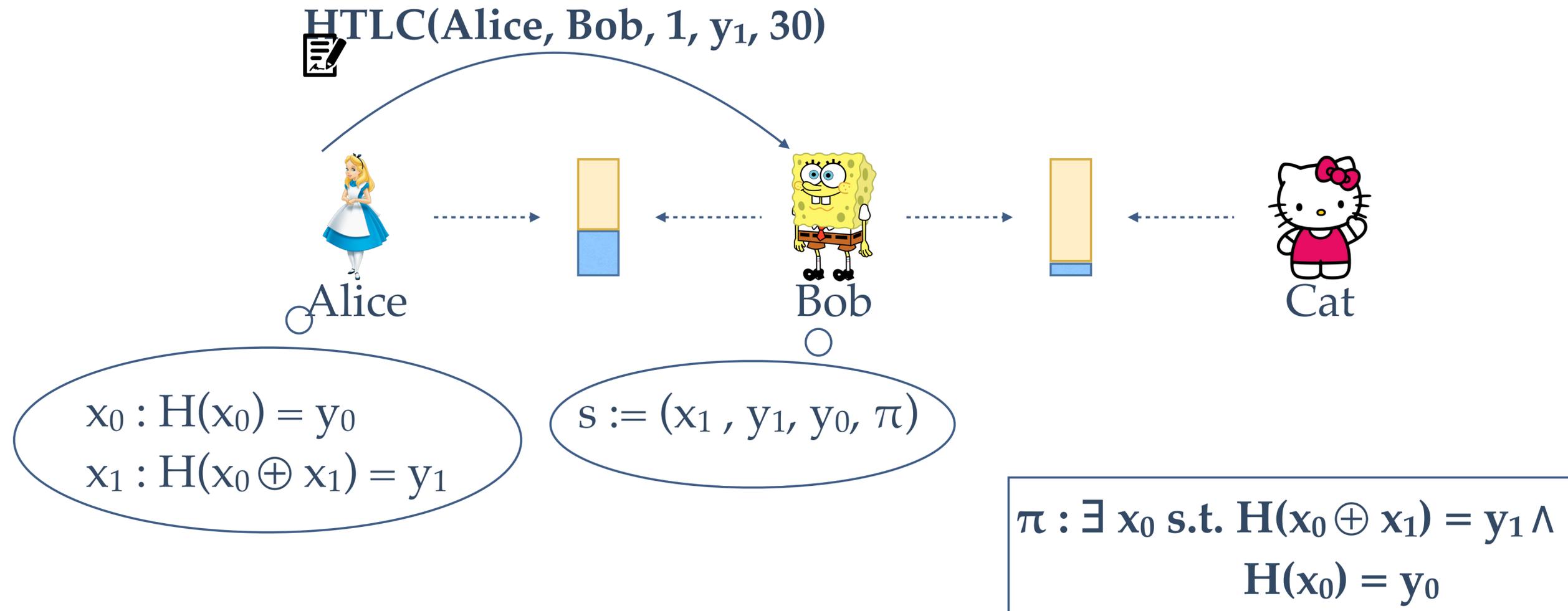
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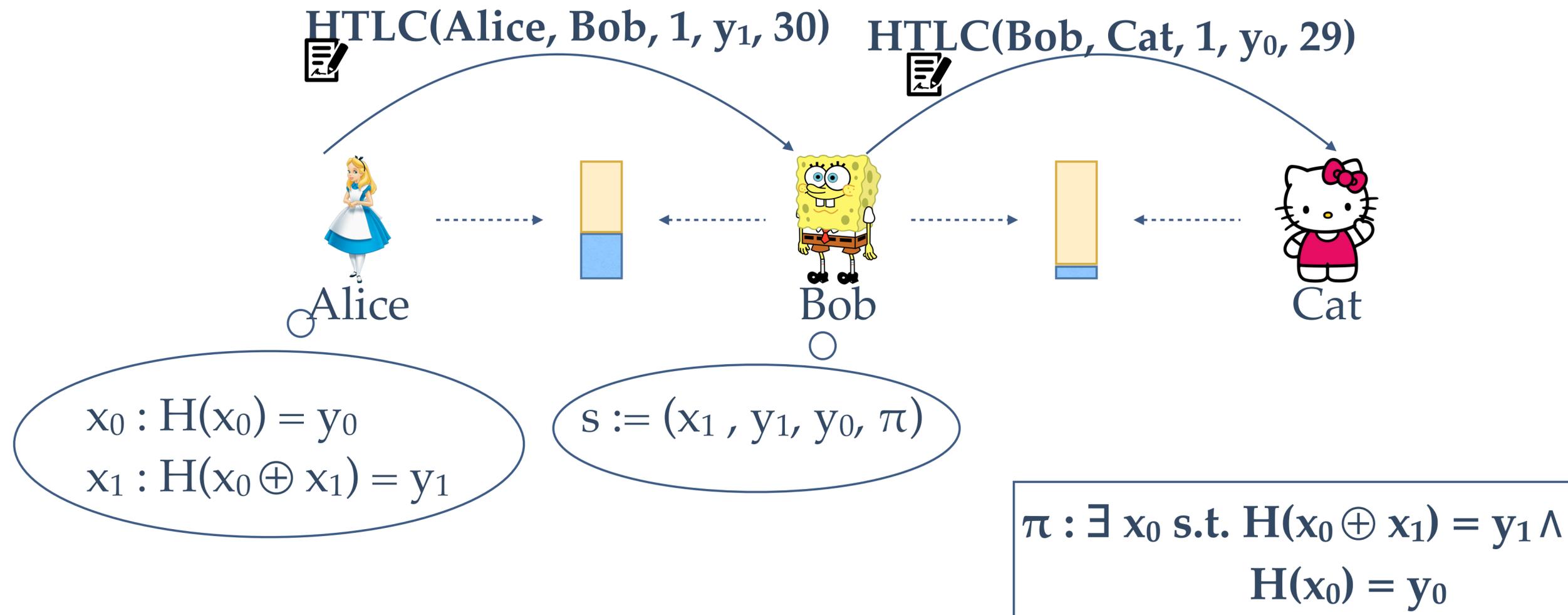
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# Multi-hop HTLC

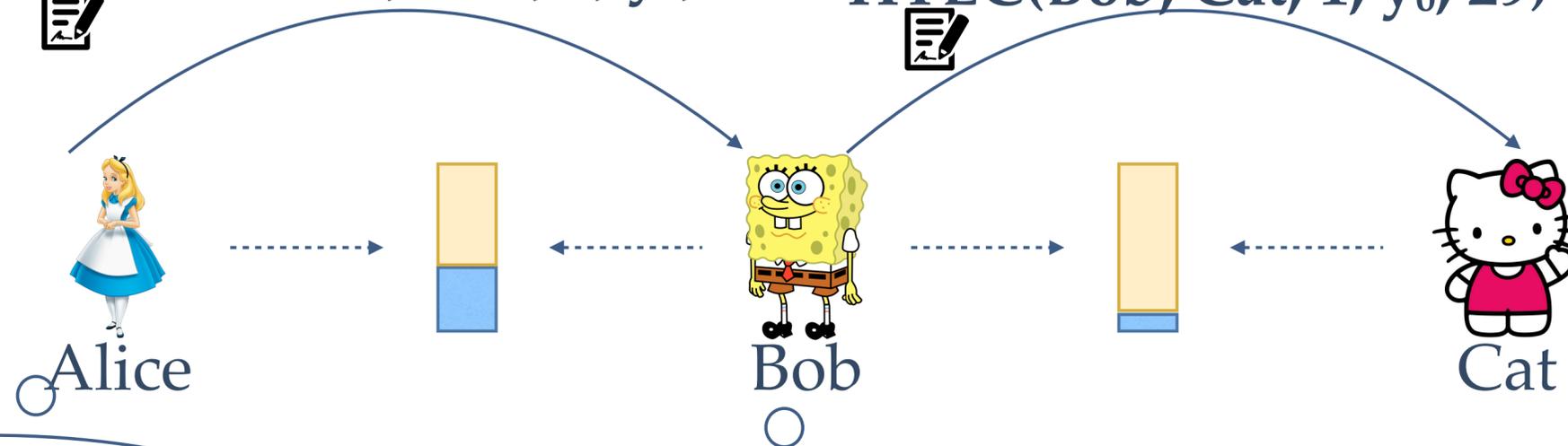
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# Multi-hop HTLC

HTLC(Alice, Bob, 1,  $r_1$ , 30) HTLC(Bob, Cat, 1,  $r_0$ , 29)

$\approx$  HTLC(Alice, Bob, 1,  $y_1$ , 30) HTLC(Bob, Cat, 1,  $y_0$ , 29)

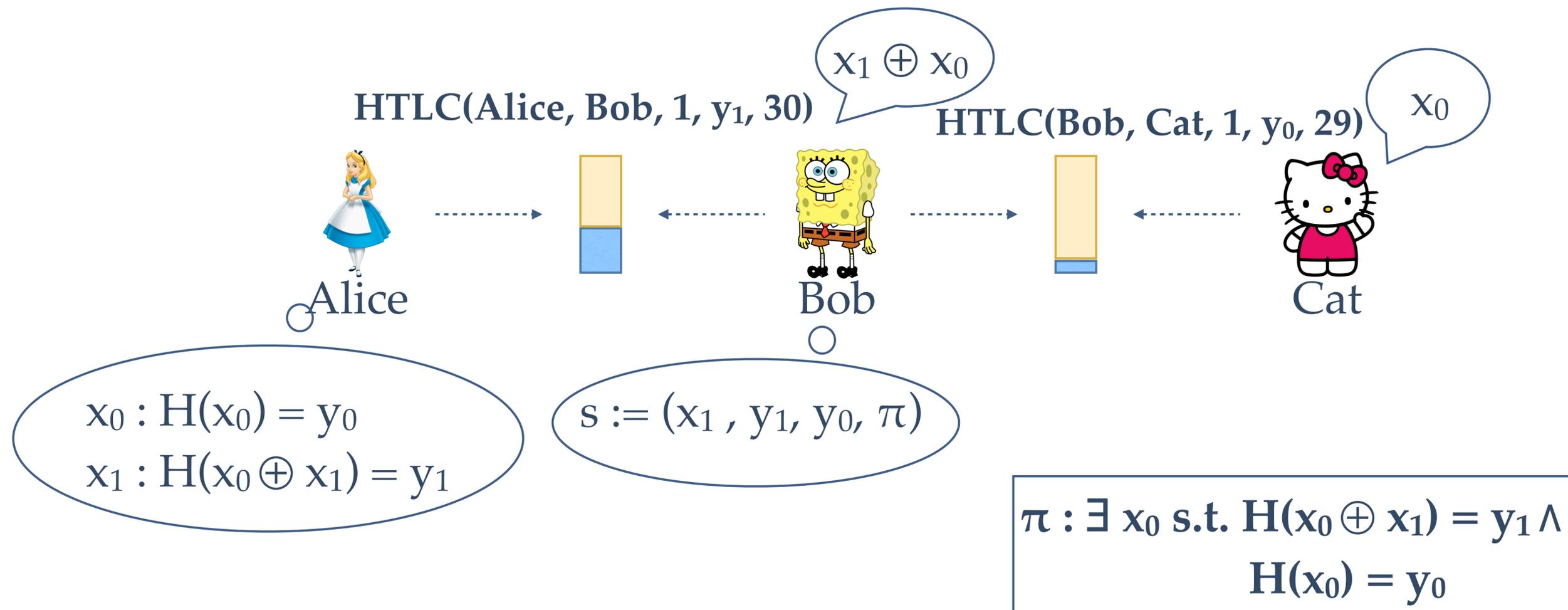


$$\begin{aligned}
 x_0 &: H(x_0) = y_0 \\
 x_1 &: H(x_0 \oplus x_1) = y_1
 \end{aligned}$$

$$s := (x_1, y_1, y_0, \pi)$$

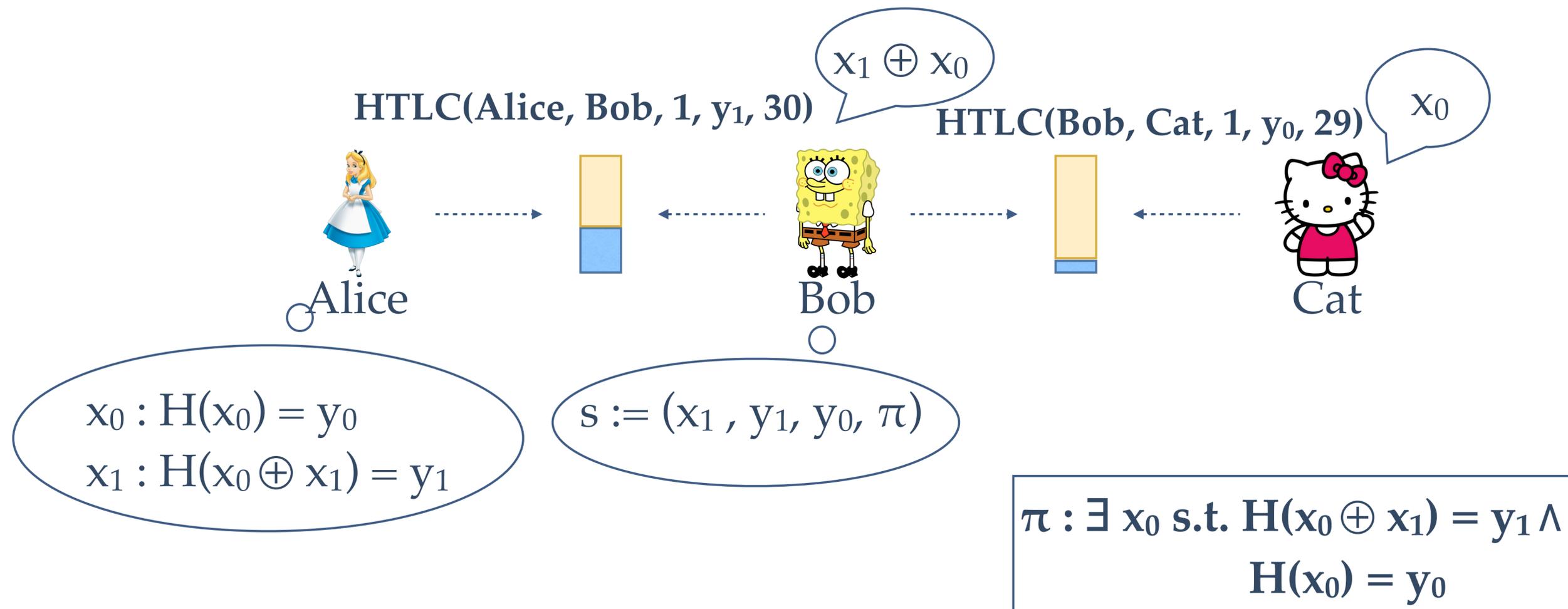
$$\begin{aligned}
 \pi &: \exists x_0 \text{ s.t. } H(x_0 \oplus x_1) = y_1 \wedge \\
 &H(x_0) = y_0
 \end{aligned}$$

# Multi-hop HTLC



# Multi-hop HTLC

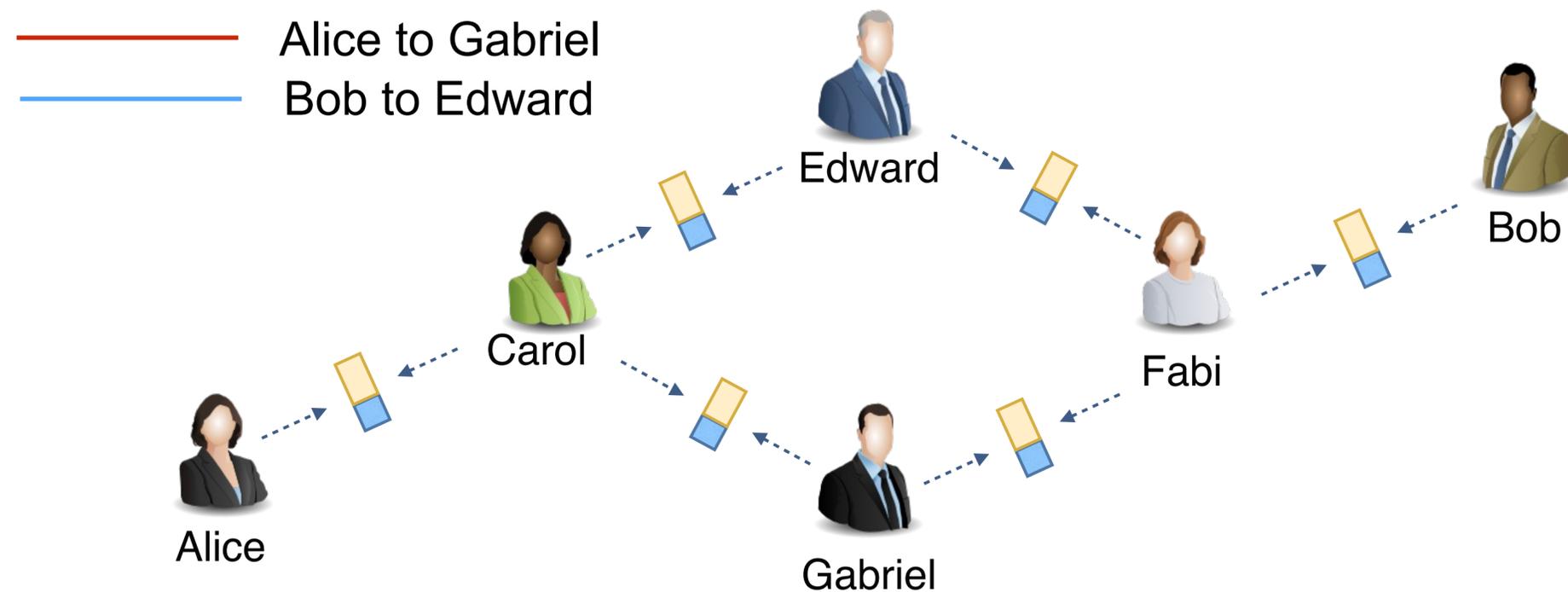
- ❖ Soundness of NIZK  $\Rightarrow$  Bob does not lose coins
- ❖ Zero-knowledge of NIZK  $\Rightarrow$  Bob does not steal coins



# Concurrency in PCNs

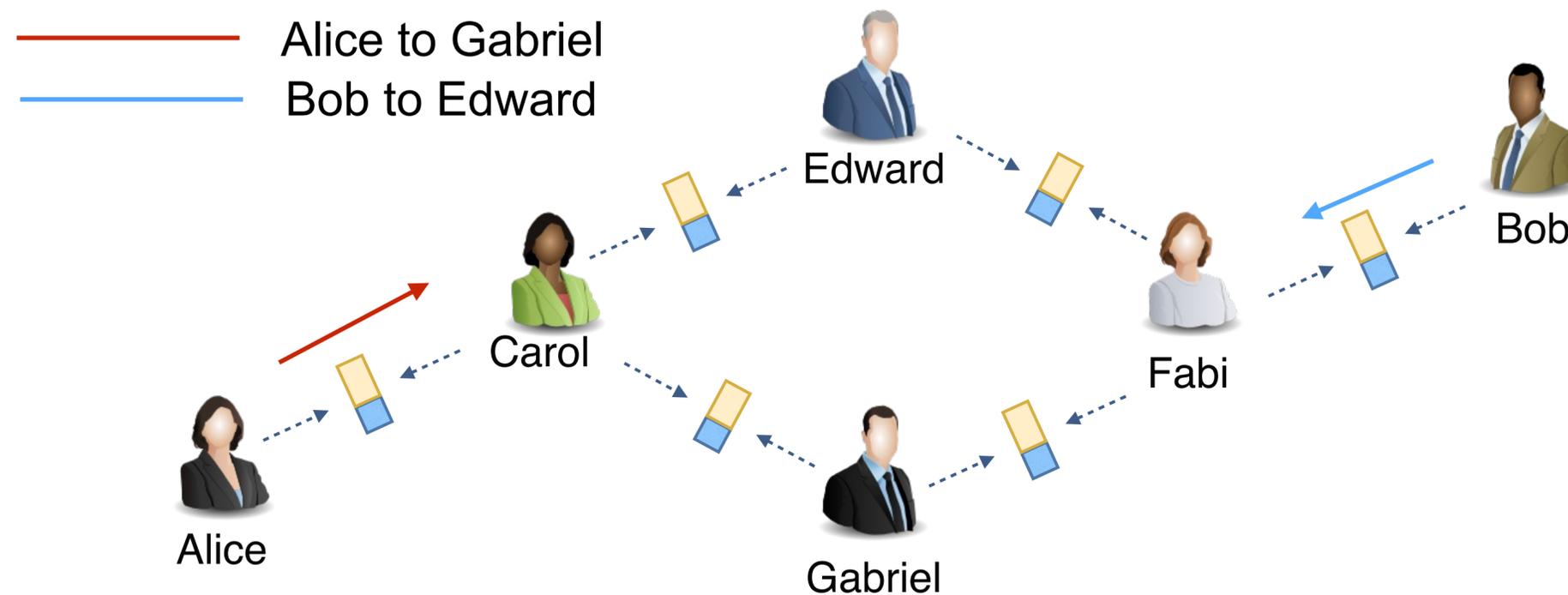
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- ❖ Concurrent on-chain payments can be easily ordered by miners
- ❖ No user has a complete view of off-chain concurrent payments in a P2P network
- ❖ A blocking solution can lead to deadlocks



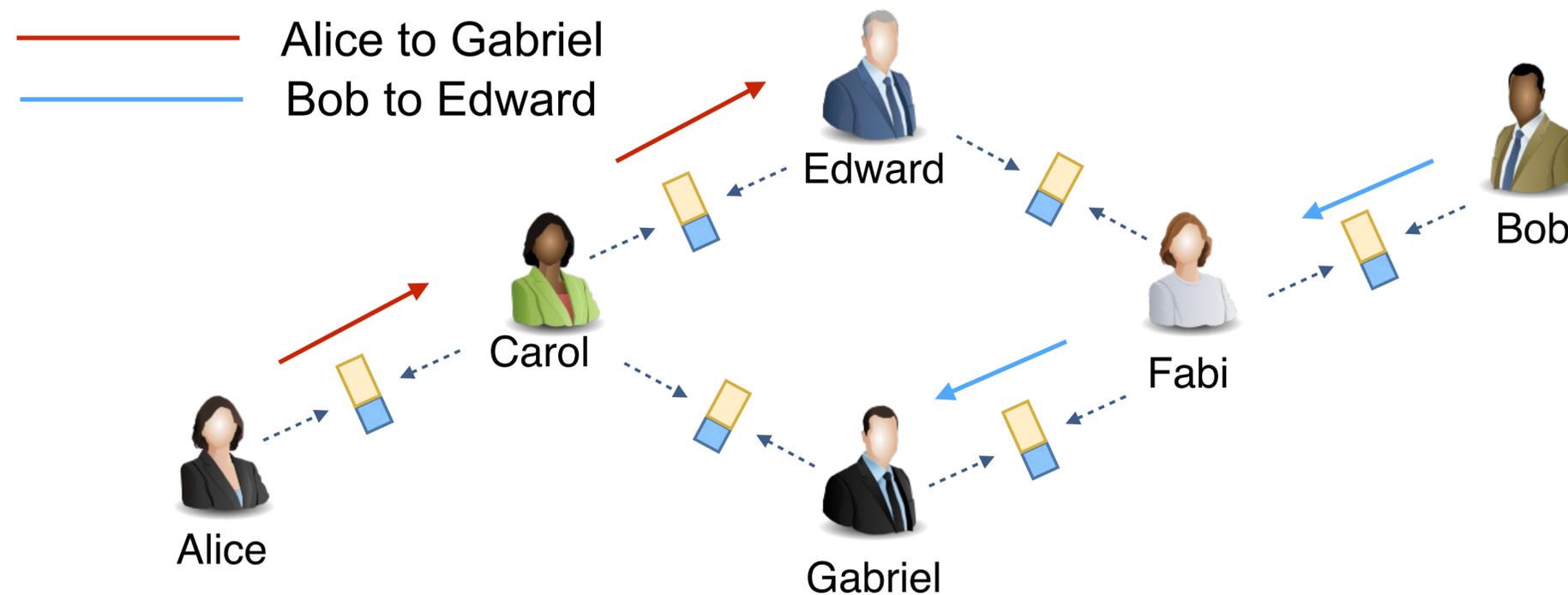
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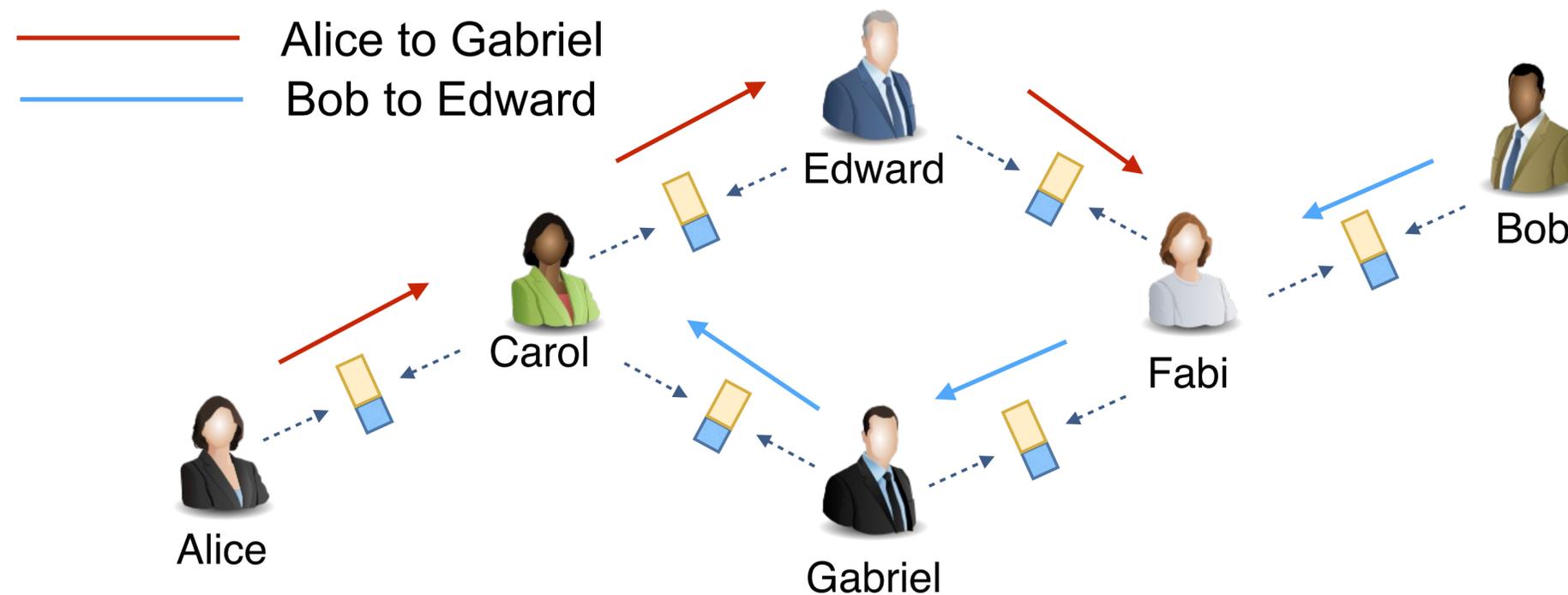
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# Concurrency in PCNs

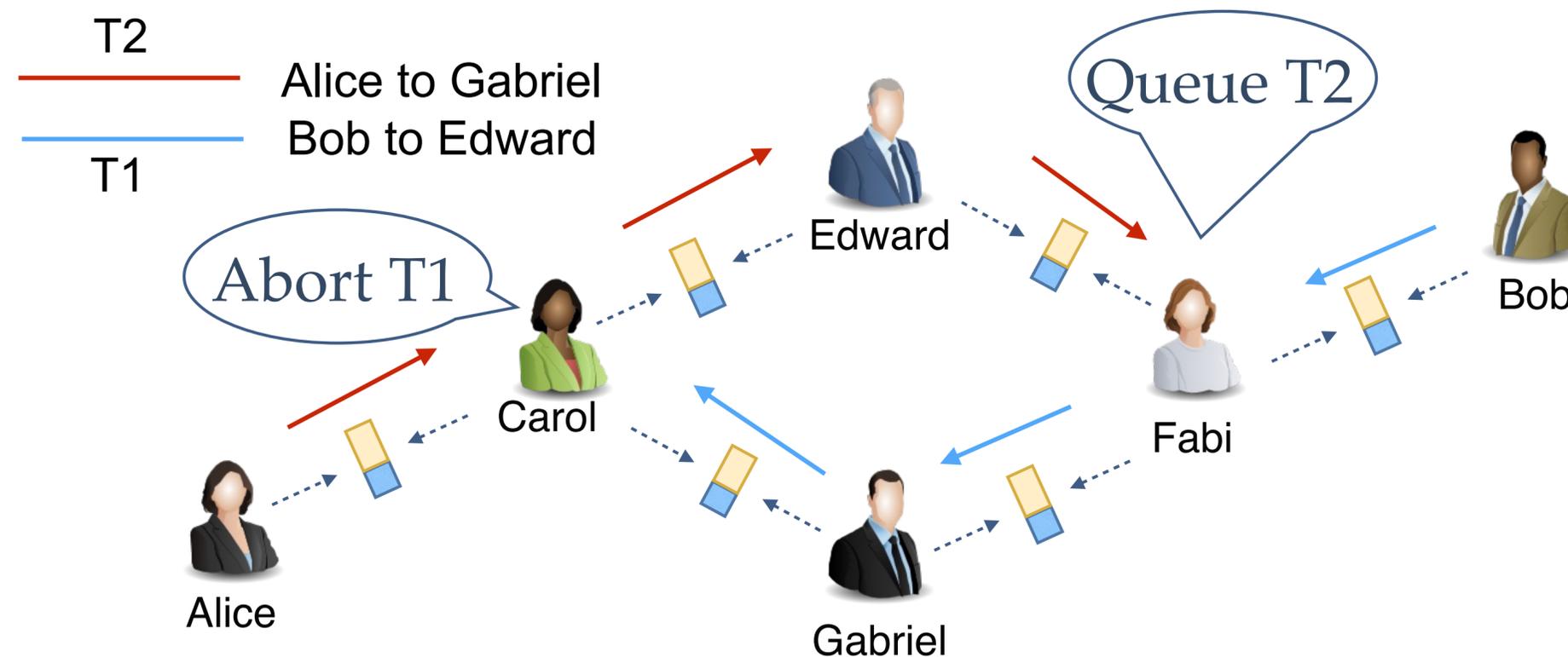
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# Concurrency in PCNs: Our Solution

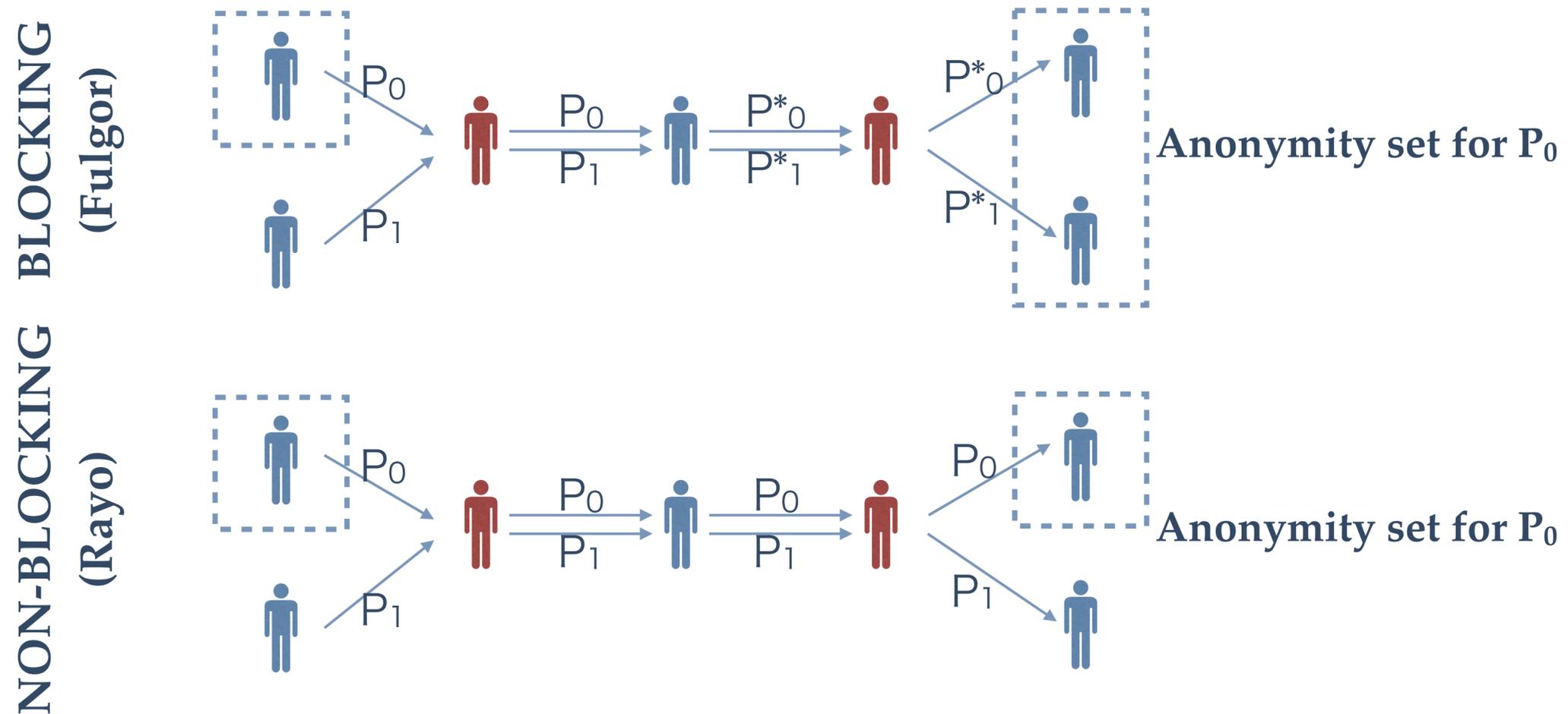
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- ❖ A non-blocking solution (Rayo): at least one payment finishes
- ❖ Main idea: Use global transaction identifiers



# Concurrency vs Privacy Tradeoff

- ❖ Global identifiers leak transaction ID to intermediate users
- ❖ Non-blocking solutions cannot achieve strong privacy



# Implementation and Performance

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- ❖ Running time of our solution largely dominated by NIZK
  - ❖ Creating a proof requires 309 ms. Proof verification requires 130 ms
  - ❖ Proof size: 1.65MB
- ❖ 5-hop payment:
  - ❖ **Non-private (LN): 609 ms**
  - ❖ **Private: 1929 ms and ~ 5 MB (Proofs are not included in the blockchain)**

# Conclusions

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- ❖ Define the security and privacy properties of interest in PCN
- ❖ Inherent tradeoff between concurrency and privacy
- ❖ Fulgor and Rayo: two approaches for concurrency and privacy
- ❖ Our solutions are efficient, compatible with Bitcoin script and without storage overhead in the blockchain

*Thank you for your attention!*

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